# Fast parallel bio-molecular solutions: the set-basis problem 

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#### Abstract

In the paper, it is demonstrated how to apply sticker in the sticker-based model for constructing solution space of DNA for the setbasis problem and how to apply DNA operations in the Adleman-Lipton model to solve that problem from solution space of sticker. Furthermore, this work shows the ability of DNA-based computing for resolving the NP-complete problems.


Keywords: biological parallel computing; DNA-based supercomputer; NP-complete problem; set-basis problem.

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#### Abstract

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## 1 Introduction

Nowadays, producing roughly $10^{18}$ DNA strands, that too in a test tube, through advances in molecular biology is possible (Sinden, 1994). Basic biological operations can be applied to simultaneously operate $10^{18}$ bits of information. This is to say that there are $10^{18}$ data processors to be parallelly executed. Hence, it is very clear that biological computing can provide a very similar parallelism for dealing with the problem in the real world.

Adleman (1994) wrote the first paper that DNA strands could be used to deal with solutions for an instance of the NP-complete Hamiltonian path problem (HPP). Lipton (1995) wrote the second paper that demonstrated that the Adleman techniques could be used to solve the NP-complete satisfiability (SAT) problem (the first NP-complete problem). Adleman and his coauthors (Roweis et al., 1999) proposed sticker for enhancing the Adleman-Lipton model.

In this paper, we use a sticker in the sticker based model for constructing a solution space of DNA for the setbasis problem. Simultaneously, we also apply DNA operations in the Adleman-Lipton model to develop a DNA algorithm. It is shown from the main result of the proposed DNA algorithm that the setbasis problem is resolved with biological operations in the Adleman-Lipton model from the solution space of the sticker. Furthermore, this work shows the ability of DNAbased computing for resolving the NP-complete problems.

The rest of this paper is organised as follows. In Section 2, the Adleman-Lipton model is introduced in detail and the comparison of the model with other models is given. Section 3 introduces a DNA algorithm for solving the setbasis problem from the solution space of the sticker in the Adleman-Lipton model. In Section 4, the experimental result of simulated DNA computing is discussed. Conclusions are drawn in Section 5.

## 2 DNA model of computation

In Subsection 2.1, a summary of DNA structure is given and the Adleman-Lipton model is described in detail. In Subsection 2.2, a comparison of the Adleman-Lipton model with other models is given.

### 2.1 The Adleman-Lipton model

A DNA (DeoxyriboNucleic Acid) is the molecule that plays the main role in DNA based computing (Paun et al., 1998). In the biochemical world of large and small molecules, polymers, and monomers, DNA is a polymer, which is strung together from monomers called deoxyribonucleotides. The monomers used for the construction of DNA are deoxyribonucleotides. Each deoxyribonucleotide contains three components: a sugar, a phosphate group, and a nitrogenous base. The sugar has five carbon atoms - for the sake of reference, there is a fixed numbering for them. Because the base also has carbons, to avoid confusion, the carbons of the sugar are numbered from $1^{\prime}$ to $5^{\prime}$ (rather than from 1 to 5 ). The phosphate group is attached to the $5^{\prime}$ carbon, and the base is attached to the $1^{\prime}$ carbon. Within the sugar structure there is a hydroxyl group attached to the $3^{\prime}$ carbon.

Distinct nucleotides are detected only with their bases, which come in two sorts: purines and pyrimidines (Sinden, 1994; Paun et al., 1998). Purines include adenine and guanine, abbreviated $A$ and $G$. Pyrimidines contain cytosine and thymine, abbreviated $C$ and $T$. Because nucleotides are only distinguished from their bases, they are simply represented as $A, G, C$, or $T$ nucleotides, depending upon the sort of base that they have. The structure of a nucleotide, cited from (Paun et al., 1998), is illustrated (in a very simplified way) in Figure 1. In Figure 1, B is one of the four possible bases $(A, G, C$, or $T), P$ is the phosphate group, and the rest (the 'stick') is the sugar base (with its carbons enumerated 1 ' through $5^{\prime}$ ).

Table 3 Sequences chosen to represent the two elements in $S$ in Figure 2

| Vertex | $5^{\prime} \rightarrow 3^{\prime}$ DNA sequence |
| :--- | :---: |
| $x_{2}^{0}$ | CCTACCTCTCACCTT |
| $x_{1}^{0}$ | CCACATATCCATCCC |
| $x_{2}^{1}$ | CATTACCTCTTTACT |
| $x_{1}^{1}$ | CCCATCTTTCTTAAC |

Table 4 The energy for the binding of each probe to its corresponding region on a library strand

|  | Enthalpy energy <br> (H) | Entropy energy <br> $(\mathrm{S})$ | Free energy (G) |
| :--- | :---: | :---: | :---: |
| $x_{2}^{0}$ | 109.3 | 278.4 | 26.2 |
| $x_{1}^{0}$ | 110.9 | 278 | 28 |
| $x_{2}^{1}$ | 106.7 | 279.7 | 23 |
| $x_{1}^{1}$ | 112.1 | 288.8 | 25.8 |

The program simulates a mix and split combinatorial synthesis technique (Cukras et al., 1998) to synthesise the library strand to every possible subset. Those library strands are shown in Table 5 and, correspondingly, represent four possible subsets: $\varphi,\{1\},\{2\}$, and $\{1,2\}$. The program is also applied to figure out the average and standard deviation for the enthalpy, entropy and free energy over all probe/library strand interactions. The energy is shown in Table 6 . The standard deviation for delta $G$ is small because this is partially enforced by the constraint that there are 4,5 , or 6 Gs (the seventh constraint in Subsection 3.2) in the probe sequences.

Table 5 DNA sequences chosen represent all possible subsets

| $\begin{aligned} & 5^{\prime}-\text { CCT ACCTCTC ACCTTCC AC AT ATCC ATCCC - } 3^{\prime} \\ & 3^{\prime}-\text { GGATGGAGAGTGGAAGGTGTATAGGTAGGG }-5^{\prime} \\ & 5^{\prime}-\text { CCT ACCTCTC ACCTTCCC ATCTTTCTT AAC - } 3^{\prime} \\ & 3^{\prime}-\text { GGATGGAGAGTGGAAGGGTAGAAAGAATTG - } 5^{\prime} \\ & 5^{\prime}-\text { CATTACCTCTTTACTCC AC AT ATCC ATCCC - } 3^{\prime} \\ & 3^{\prime}-\text { GT AATGGAGAAATGAGGTGT AT AGGT AGGG - } 5^{\prime} \\ & 5^{\prime}-\text { CATTACCTCTTTACTCCC ATCTTTCTT AAC - } 3^{\prime} \\ & 3^{\prime}-\text { GTAATGGAGAAATGAGGGTAGAAAGAATTG }-5^{\prime} \\ & \hline \end{aligned}$ |
| :---: |
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Table 6 The energy over all probe/library strand interactions

|  | Enthalpy <br> energy $(H)$ | Entropy <br> energy $(S)$ | Free energy <br> $(G)$ |
| :--- | :---: | :---: | :---: |
| Average | 109.75 | 281.225 | 25.75 |
| Standard deviation | 2.02161 | 4.41807 | 1.79095 |

The Adleman program was employed for computing the distribution of the types of potential mishybridisations. The distribution of the types of potential mishybridisations is the absolute frequency of a probestrand match of length $k$ from 0 to the bit length 15 (for DNA sequences) where probes are not supposed to match the strands. The distribution was, subsequently, $52,102,85,95,105,109,77,36,13,12,2,0$,
$0,0,0$ and 0 . It is pointed out from the last five zeros that there are 0 occurrences where a probe matches a strand at $11,12,13,14$ or 15 places. This shows that the third constraint in subsection 3.2 has been satisfied. It is very clear that the number of matches peaks at five (109). That is to say that there are 109 occurrences where a probe matches a strand at five places.

It is indicated from the execution of Step 2 of simulation that the result generated by Step 2 was shown in Table 7. The goal of Step 3 is to find a set-basis from the result generated by Step 2. Hence, Step 3(a) of simulation, the set-basis was shown in Table 8. That is to say that the answer of the set-basis problem for the finite set $S$ and the collection $C$ in Figure 2 is $\{\{1\},\{2\}\}$.

Table 7 DNA sequences generated by Step 2 represent legal subsets

5' - CCT ACCTCTC ACCTTCCC ATCTTTCTT AAC - 3'
$5^{\prime}$ - CATTACCTCTTTACTCC AC AT ATCC ATCCC - $3^{\prime}$

Table 8 DNA sequence represents the answer of the set-basis problem for the finite set $S$ and the collection $C$ in Figure 2

$$
\begin{aligned}
& \hline 5^{\prime}-\text { ССТАССТСТСАССТТСССАТСТТТСТТААС }-3^{\prime} \\
& 5^{\prime}-\text { САТТАССТСТТТАСТССАСАТАТССАТССС }-3^{\prime} \\
& \hline
\end{aligned}
$$

## 5 Conclusions

Applying splints constructs the solution space of the DNA sequence for solving the NPcomplete problem in the Adleman-Lipton and this is the reason that hybridisation has higher probabilities for errors. Adleman and his coauthors (Roweis et al., 1999) proposed a sticker to decrease probabilities of errors in hybridisation in the Adleman-Lipton model. In the proposed algorithm, the size of the solution space of the sticker is exponential. Hence, this is the limit to which we can resolve the size of the NPcomplete problem. The main result of the proposed algorithm shows that the set-basis problem is resolved with biological operations in the Adleman-Lipton model from solution space of sticker. Furthermore, this work demonstrates the ability of DNA based computing to solve NPcomplete problems.

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